Institute of Communication Systems and Data Processing Prof. Dr.-Ing. Peter Vary

Modern Techniques for Flexible, Iterative Source-Channel Decoding

Laurent Schmalen and Peter Vary

FlexCode Public Seminar June 6, 2007





- Soft Decision Source Decoding
- The Turbo Concept of Iterative Decoding
- Generation of Extrinsic Information for the Source Decoder
- Iterative Source-Channel Decoder
- Adaptive, Flexible, Multi-Mode Iterative Source-Channel Decoding scheme





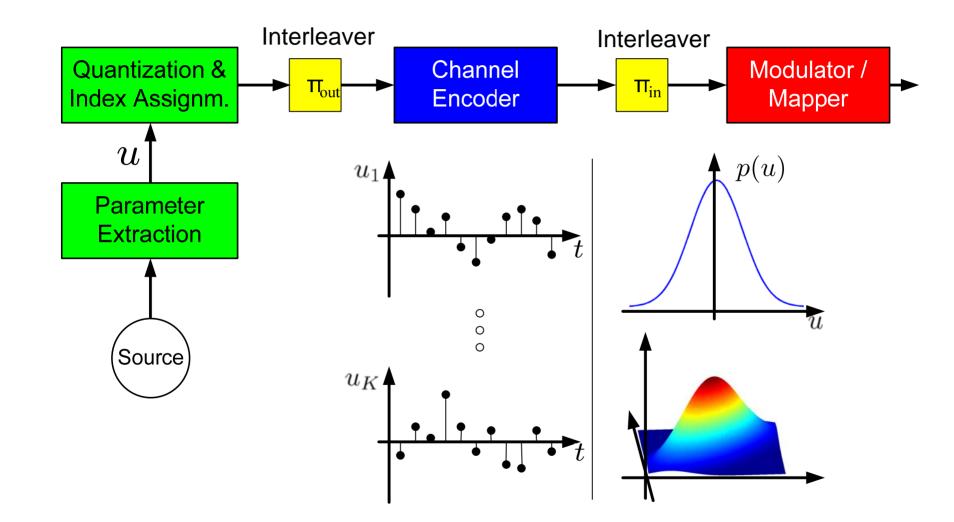
- Channel coding cannot prevent occurrence of residual bit errors in the case of adverse channel conditions leading to a severe degradation of the signal quality
- Annoying effects can be reduced or even eliminated by means of *error concealment*
- Real source coding schemes contain residual redundancy for reasons of delay, complexity and nonstationarity
- Shannon 1948:

"However, any redundancy in the source will usually help if it is utilized at the receiving point. [...] redundancy will help combat noise."

Conventional Transmission

FlerCade

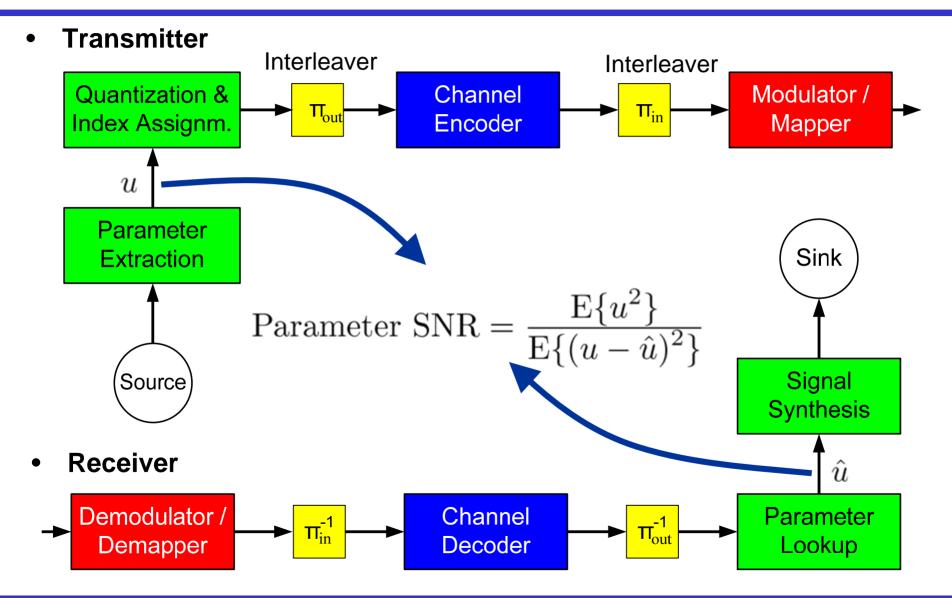


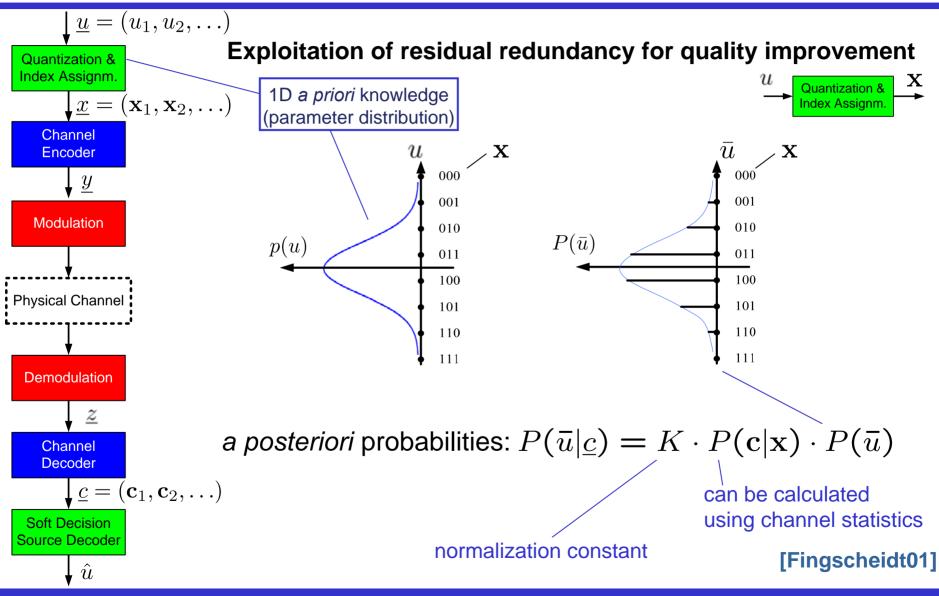


Conventional Transmission

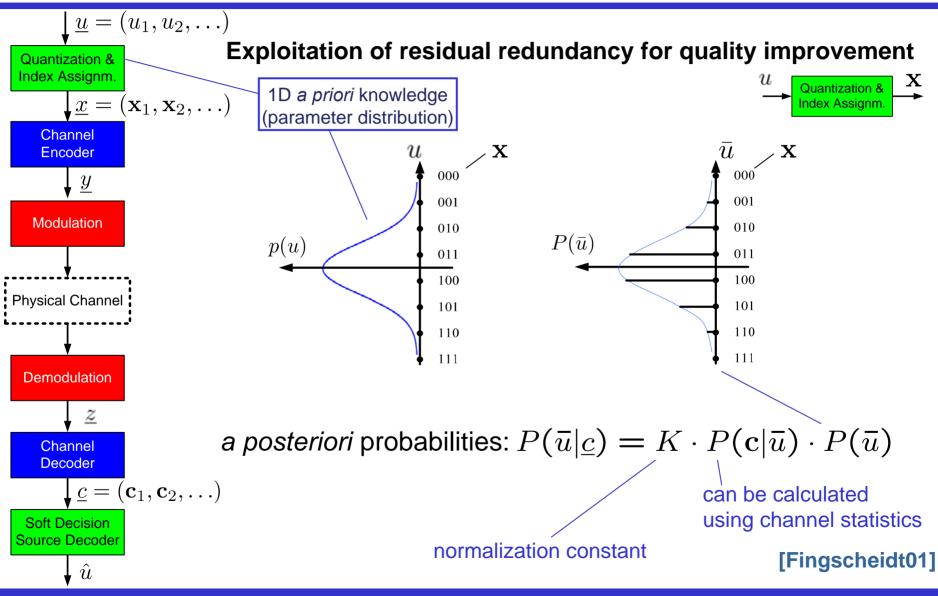
FlerCade



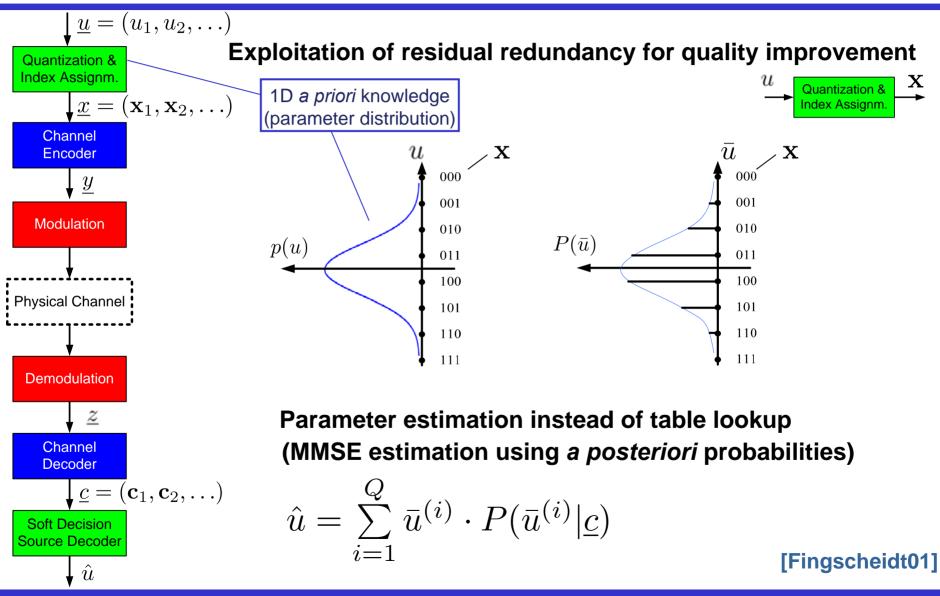




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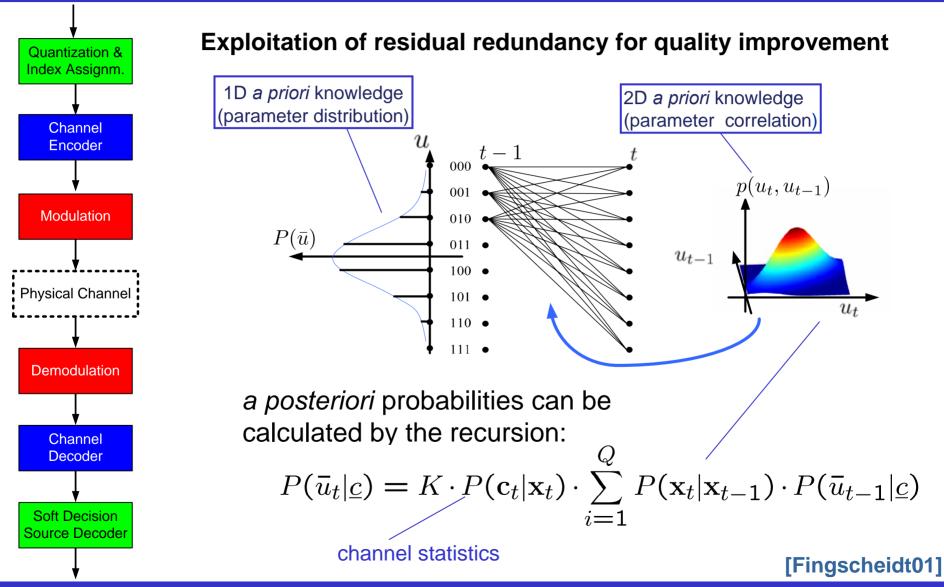


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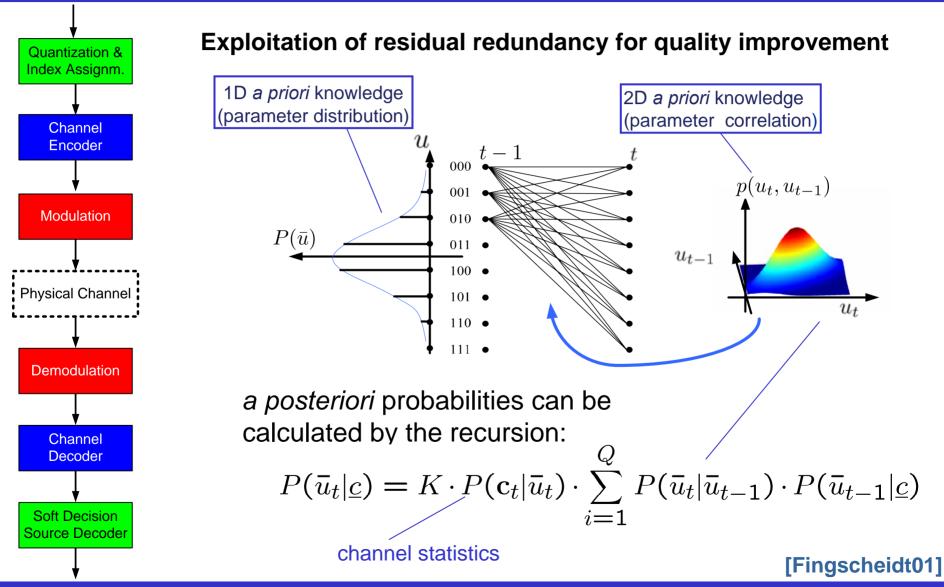
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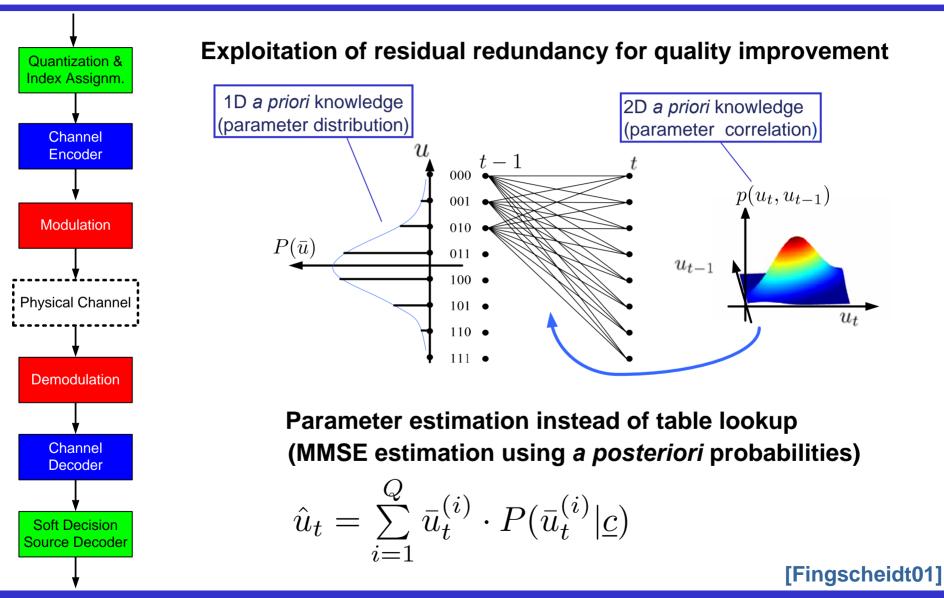
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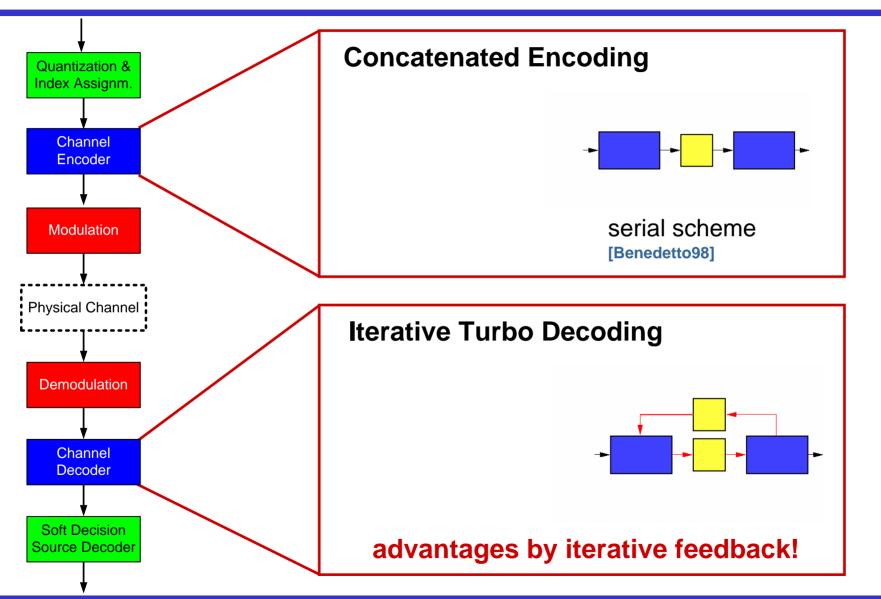


- Timeline of iterative decoding techniques
 - LDPC codes introduced in 1963[Gallager63] but forgotten due to the relatively high complexity at that time
 - Turbo codes invented in 1993[Berrou93]. Allow near-Shannon limit decoding with moderate complexity
 - LDPC codes rediscovered in 1998[MacKay98]. Decoding is also performed iteratively using belief propagation
 - Extension of the iterative decoding to other receiver components, e.g.
 - equalization (Turbo Equalization) [Douillard95]
 - modulation (BICM-ID) [Xi98]
 - multi-user detection (Turbo-MUD) [Alexander98]
 - source decoding (ISCD) [Adrat01], [Goertz01], [Guyader01]



Turbo Codes, Concept



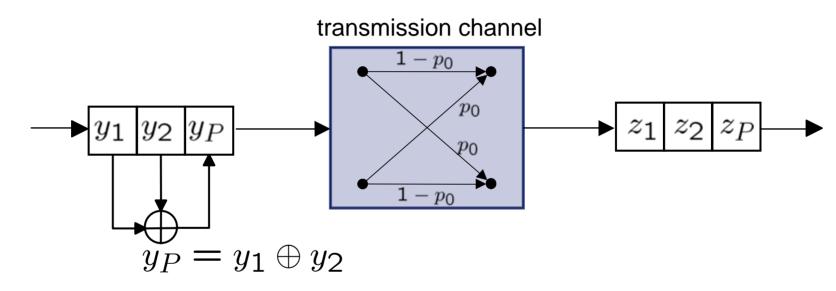






"Information from neighboring bit positions"

• Example: Parity Check Code

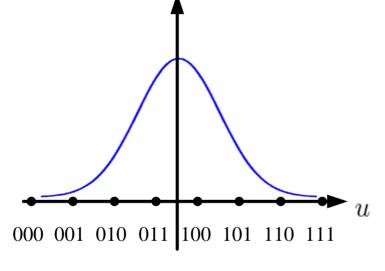


- Extrinsic information:
 - transmitter $\Rightarrow y_1 = y_2 \oplus y_P$
 - receiver $\Rightarrow \hat{z}_{e1} = z_2 \oplus z_P$

Source decoder [Adrat03] p(u)Information from above

Extrinsic Information from SDSD

Generation of extrinsic information by the soft decision



FlerCade

lacksquare

Information from channel decoder:

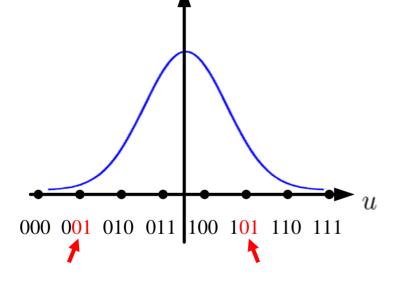
2 rightmost bits are 01



$\mathbf{A}^{p(u)}$ Information from char

Extrinsic Information from SDSD

Generation of extrinsic information by the soft decision



source decoder [Adrat03]

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lacksquare

Information from channel decoder:

2 rightmost bits are 01



 $P(x_1=0|(x_2,x_3)=(0,1)) < P(x_1=1|(x_2,x_3)=(0,1))$

Extrinsic Information from SDSD

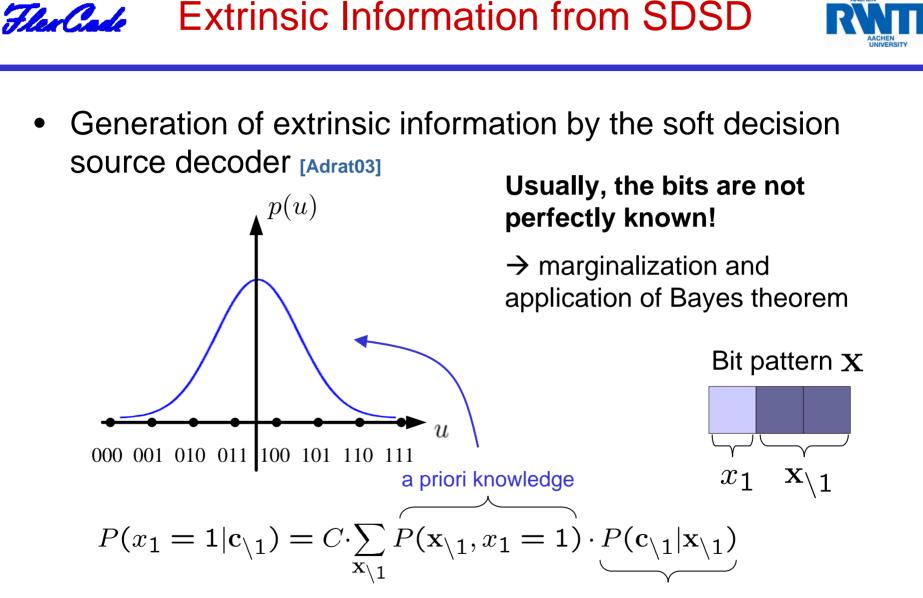
- Generation of extrinsic information by the soft decision source decoder [Adrat03]
 - p(u)0 01 1 01

Information from channel decoder:

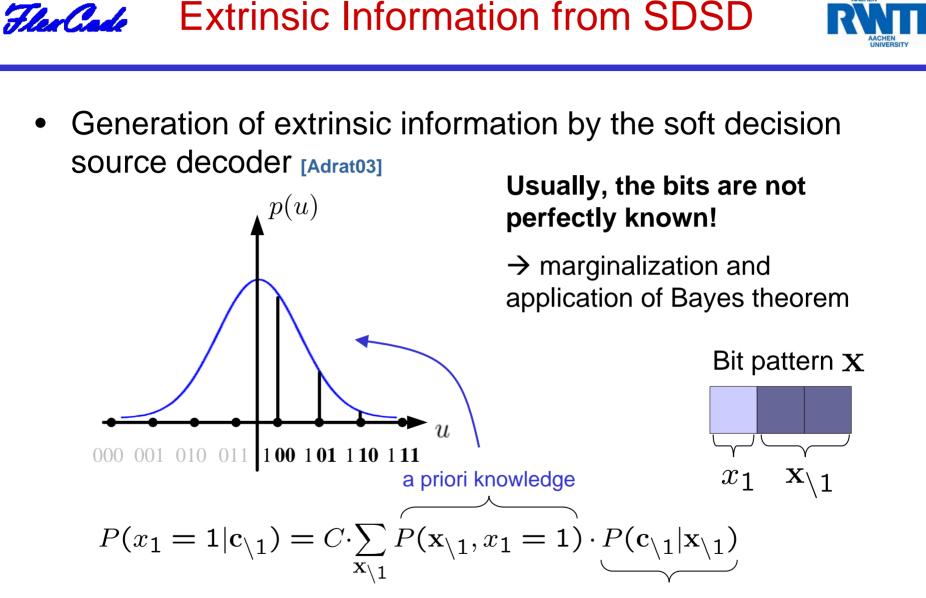
2 rightmost bits are 01

This information can be fed back to the channel decoder to improve channel decoding!





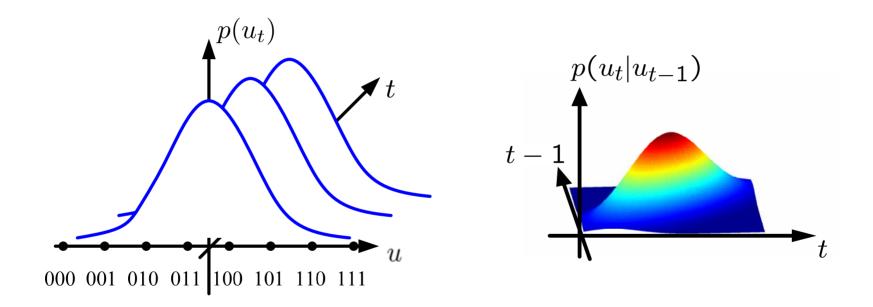
channel transition probabilities



channel transition probabilities



Correlation of the source is exploited by the SDSD

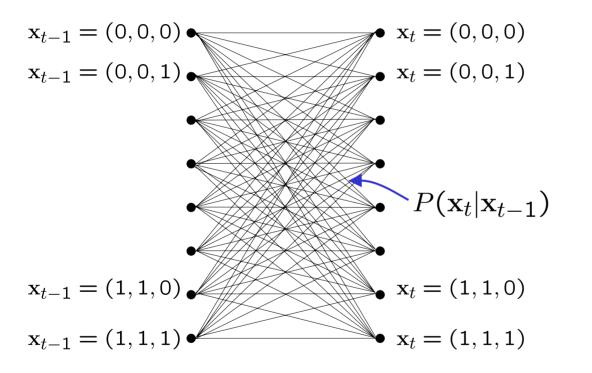


• The quantization of the source modelled as Markov process can be represented using a trellis diagram

Extrinsic Information from SDSD



• Trellis representation of the source



states correspond to quantizer reproduction levels

state transitions correspond to conditional probabilities

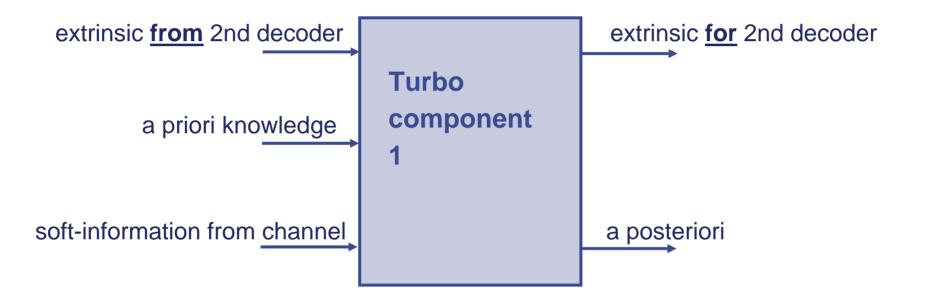
Decoding using the BCJR (MAP) algorithm [Bahl et al. 74], [Heinen 00], [Adrat 05]

• The quantization of the source modelled as Markov process can be represented using a trellis diagram

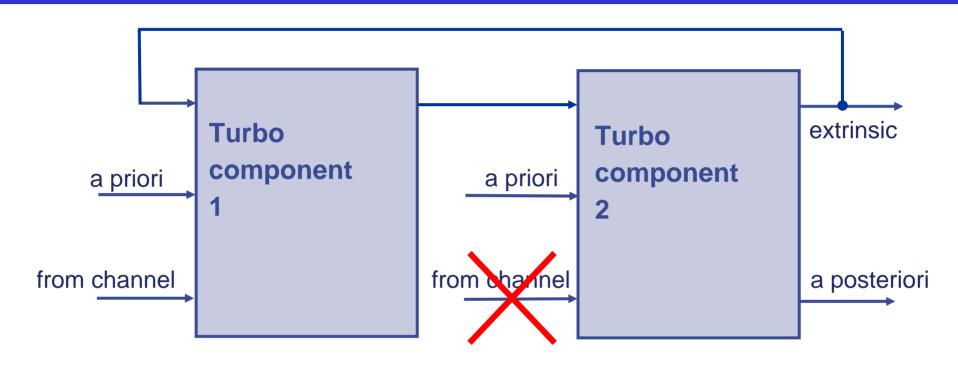




Extrinsic Information from a second decoder"



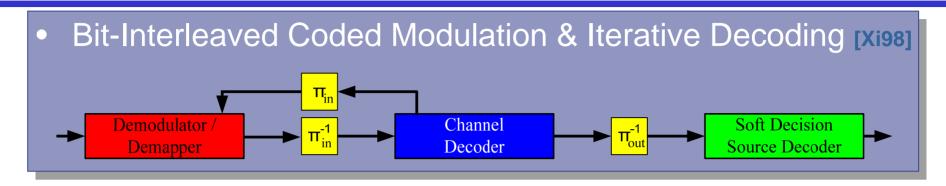
The Code Iterative Exchange of Extrinsic Information

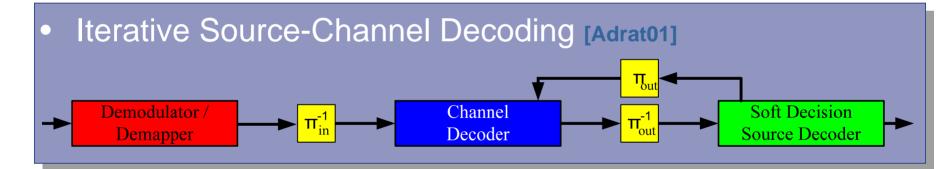


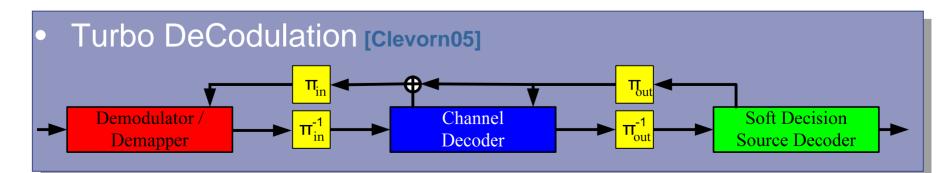
• a posteriori decision after several iterations











250 parameters/frame 13 - auto-correlation $\rho = 0.9$

15 $\mathcal{P}^{[ext{ref}]}$ AWGN / BPSK Modulation dB **ISCD:** Iterative 3 Bit Lloyd-Max guantizer **Source-Channel** - $r_{\rm c}$ = 1/2 convolutional code with 8 trellis states

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convolutional code **EXIT** optimized index assignment

first order Markov model

- recursive non-systematic

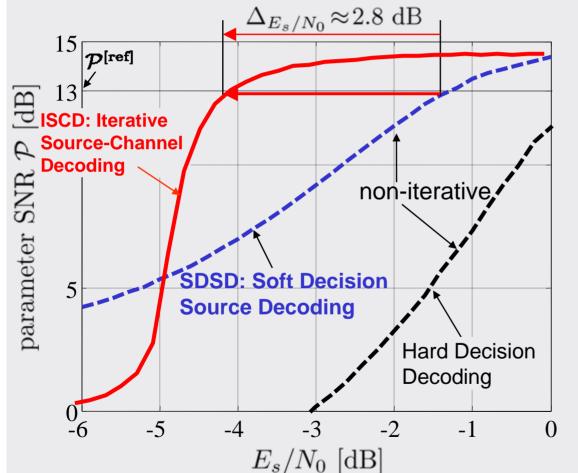
SDSD:

ISCD:

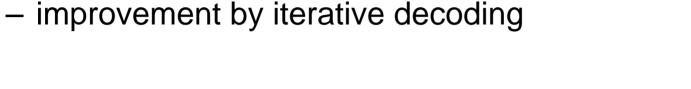
10 iterations

HarCade Iterative Source-Channel Decoding

Design constraint: Parameter SNR e.g. $\mathcal{P}^{[ref]} > 13 \text{ dB}$







Audio Examples

	Speech	Music
1 iteration		A
2 iterations		٩.
3 iterations	4	W
4 iterations	Ŷ	- C

FlerCado

Audio examples

A-law PCM with 8-bit quantization

AWGN / BPSK Modulation

- 44.1 kHz sampling rate
- 300 samples/frame
- $r_{\rm C}$ = 1 convolutional code with 8 trellis states
- redundant $r_{IA} = \frac{1}{2}$ block coded index assignment

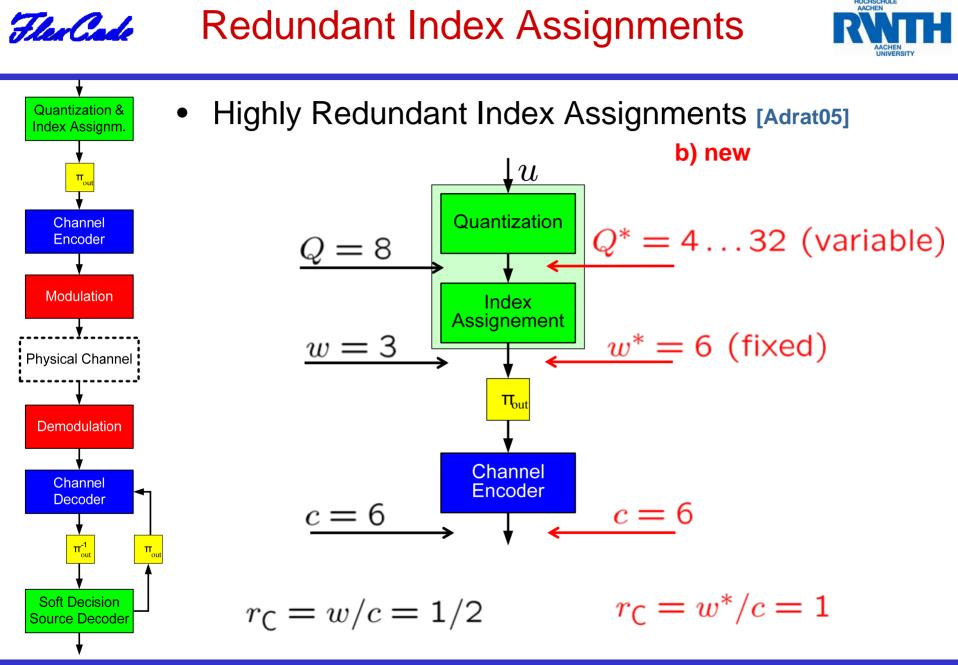
$$- E_s/N_0 = 1 \text{ dB} (\text{BER} = 5.5 \%)$$

SDSD:

 exploiting unequal parameter distribution only (zeroth order apriori knowledge)

\rightarrow further improvement by exploiting correlation

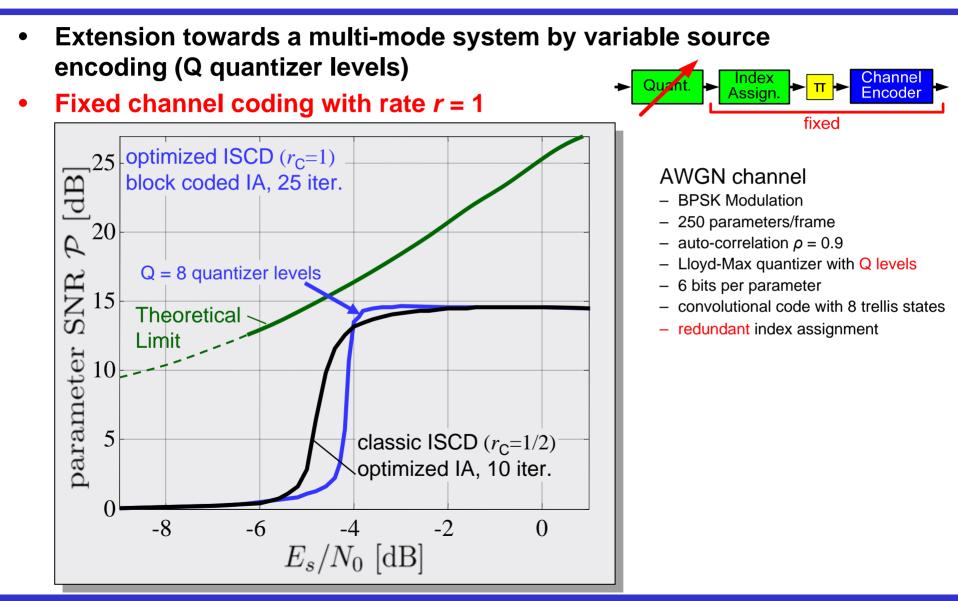
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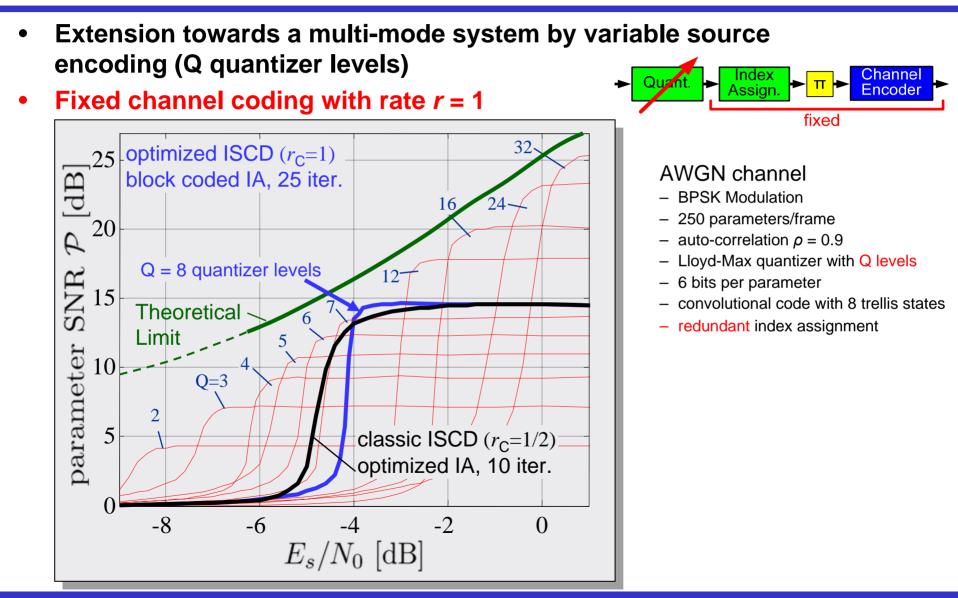
HarCade Iterative Source-Channel Decoding





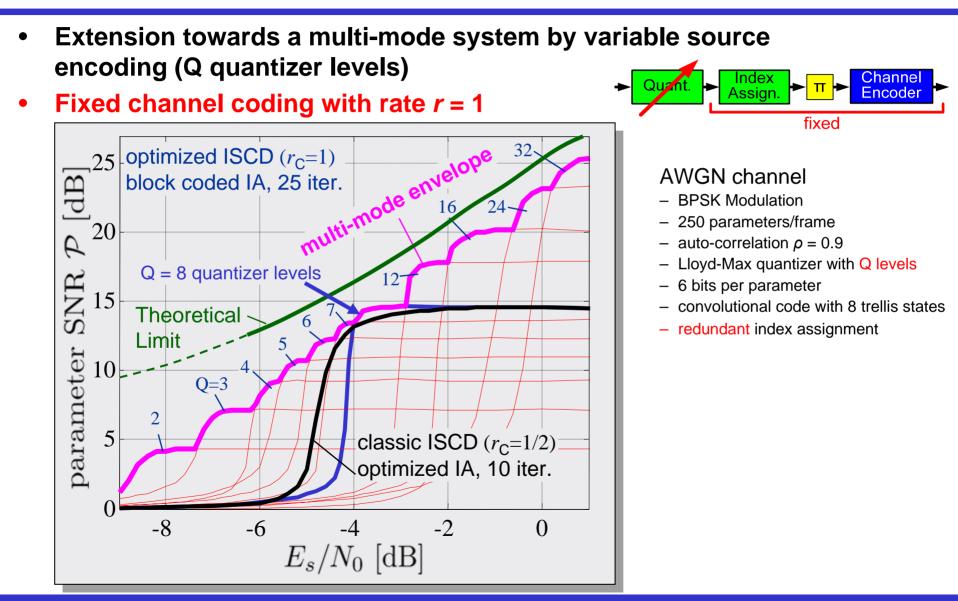
HarCade Iterative Source-Channel Decoding





HarCade Iterative Source-Channel Decoding









- Exploitation of residual source redundancy for soft decision source decoding
- Determination of extrinsic information by using residual source redundancy
- Iterative decoding concept extended to the source decoding step
 - Near capacity decoding of quantized, correlated sources
 - Adaptivity/Flexibility by using a multi-mode transmission scheme

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