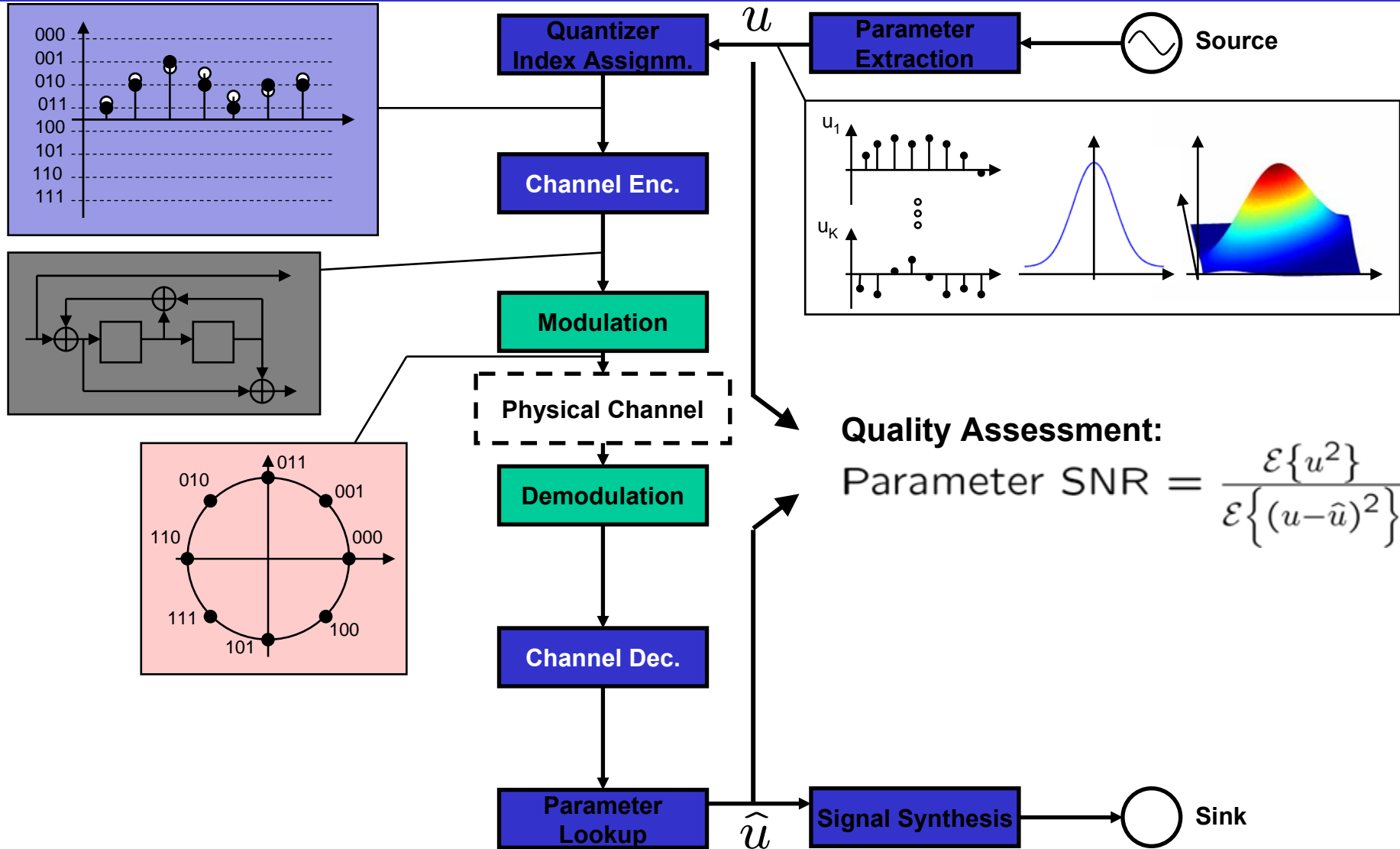


Flexible, Iterative Source-Channel Decoding

Peter Vary, Thorsten Clevorn, Laurent Schmalen
Flexcode Seminar Helsinki
November 29, 2006

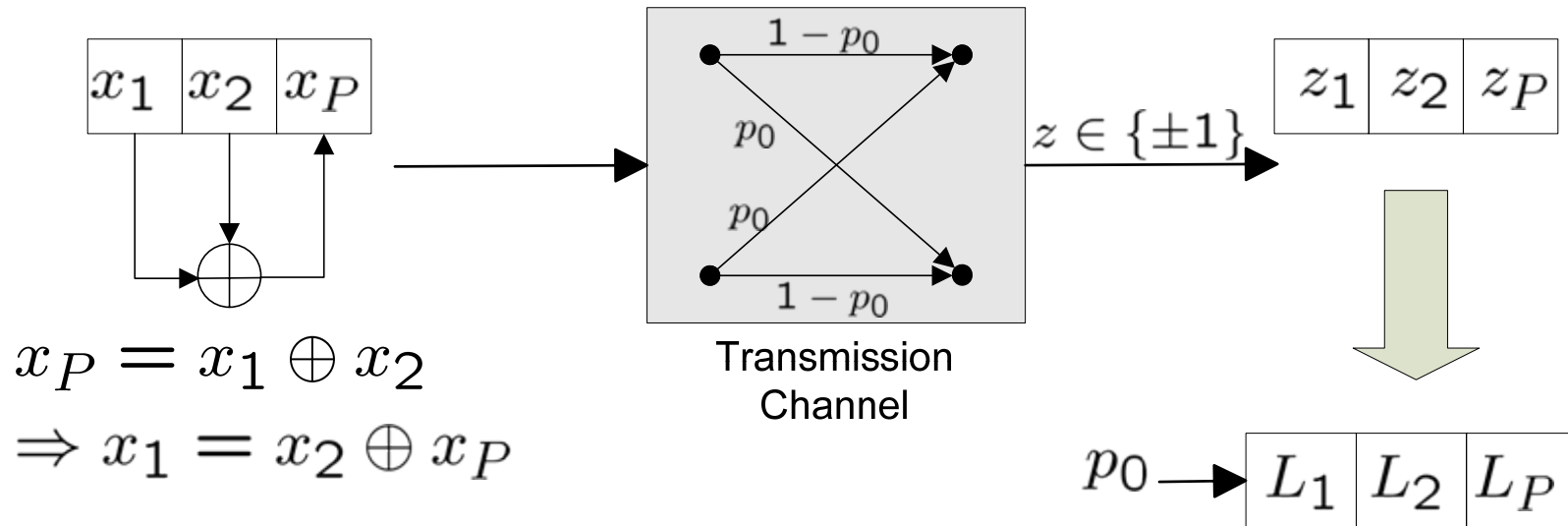
- **The Turbo Principle**
- **Iterative Source-Channel Decoding (ISCD)**
- **Multi-Mode ISCD**
- **Conclusions**



The Turbo Principle

“Information on a bit originating from other bit positions”

- Extrinsic Information is the key element in Turbo decoders
- Example: Parity Check



Reliability information

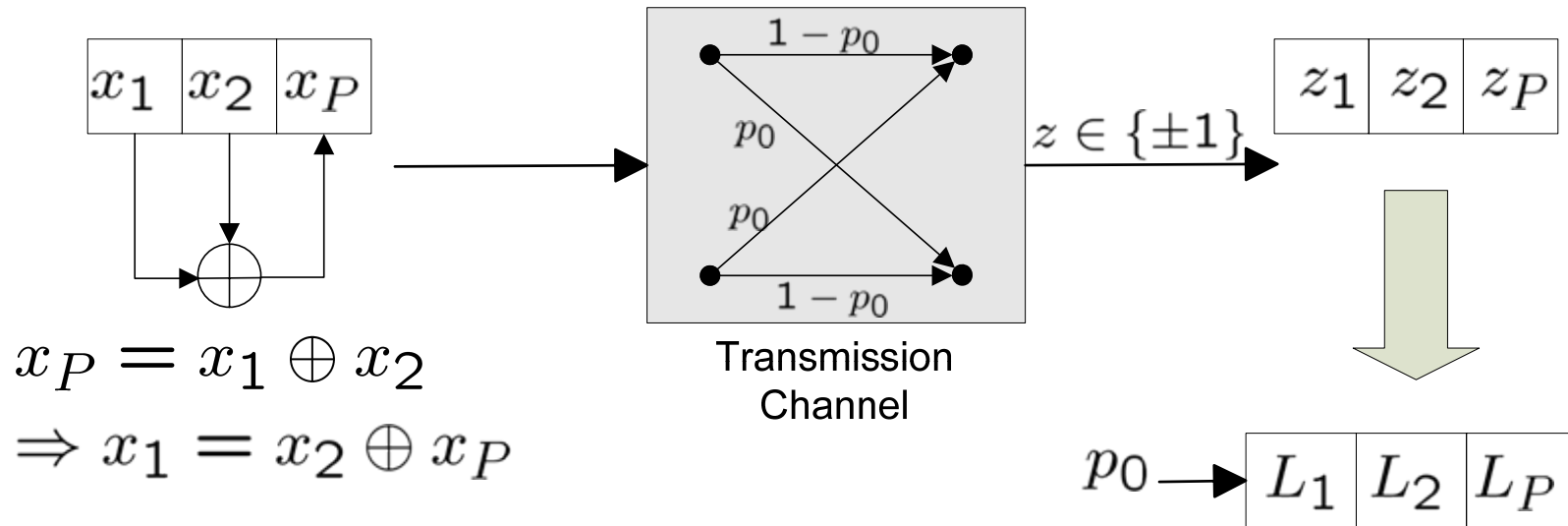
$$L_1 = z_1 \cdot \ln \frac{1-p_0}{p_0}$$

$$L_2 = z_2 \cdot \ln \frac{1-p_0}{p_0}$$

$$L_P = z_P \cdot \ln \frac{1-p_0}{p_0}$$

“Information on a bit originating from other bit positions”

- Extrinsic Information is the key element in Turbo decoders
- Example: Parity Check



Reliability information

$$L_1 = z_1 \cdot \ln \frac{1-p_0}{p_0}$$

$$L_2 = z_2 \cdot \ln \frac{1-p_0}{p_0}$$

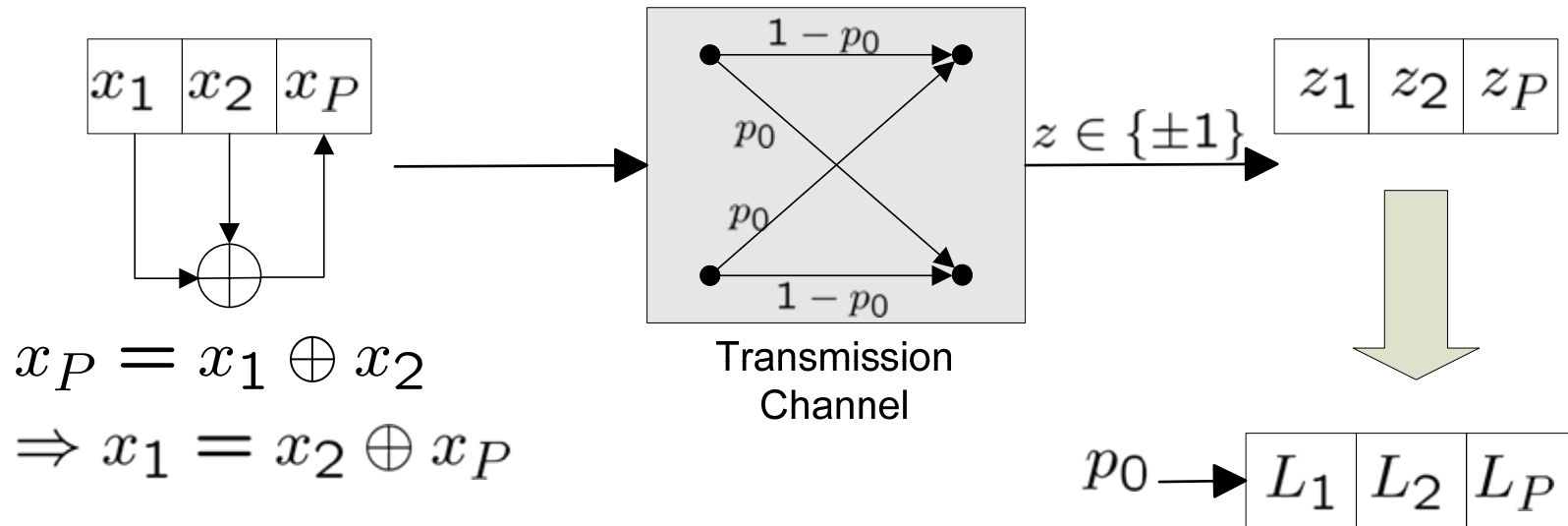
$$L_P = z_P \cdot \ln \frac{1-p_0}{p_0}$$

→ intrinsic: $L_i = L(z_i | x_i)$

$$L_i = \ln \left(\frac{P(z_i | x_i = +1)}{P(z_i | x_i = -1)} \right)$$

“Information on a bit originating from other bit positions”

- Extrinsic Information is the key element in Turbo decoders
- Example: Parity Check



Reliability information

$$L_1 = z_1 \cdot \ln \frac{1-p_0}{p_0}$$

→ intrinsic: $L_i = L(z_i|x_i)$

$$L_2 = z_2 \cdot \ln \frac{1-p_0}{p_0}$$

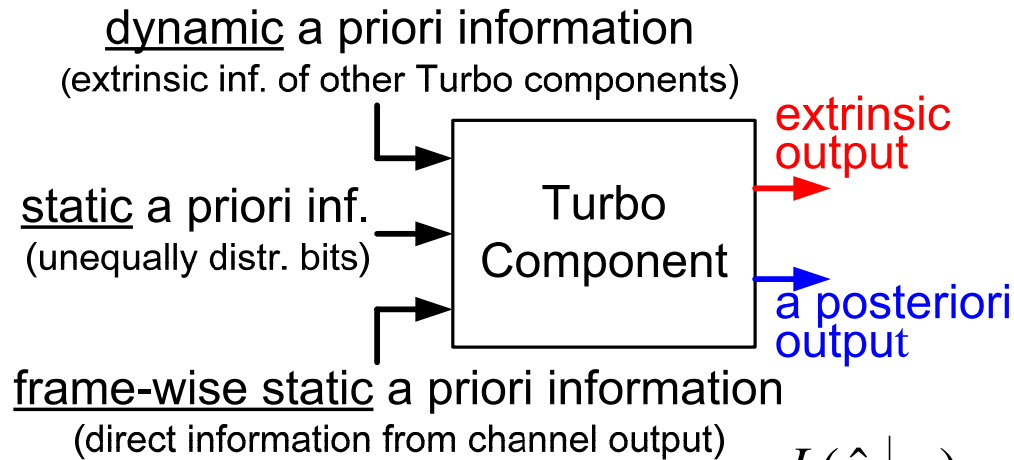
→ a posteriori $L(\hat{x}_i|z_i) = L(z_i|x_i) + L(x_i) + L_i^{[ext]}$

$$L_P = z_P \cdot \ln \frac{1-p_0}{p_0}$$

→ $\hat{x}_i = \text{sign}\{L(\hat{x}_i|z_i)\}$

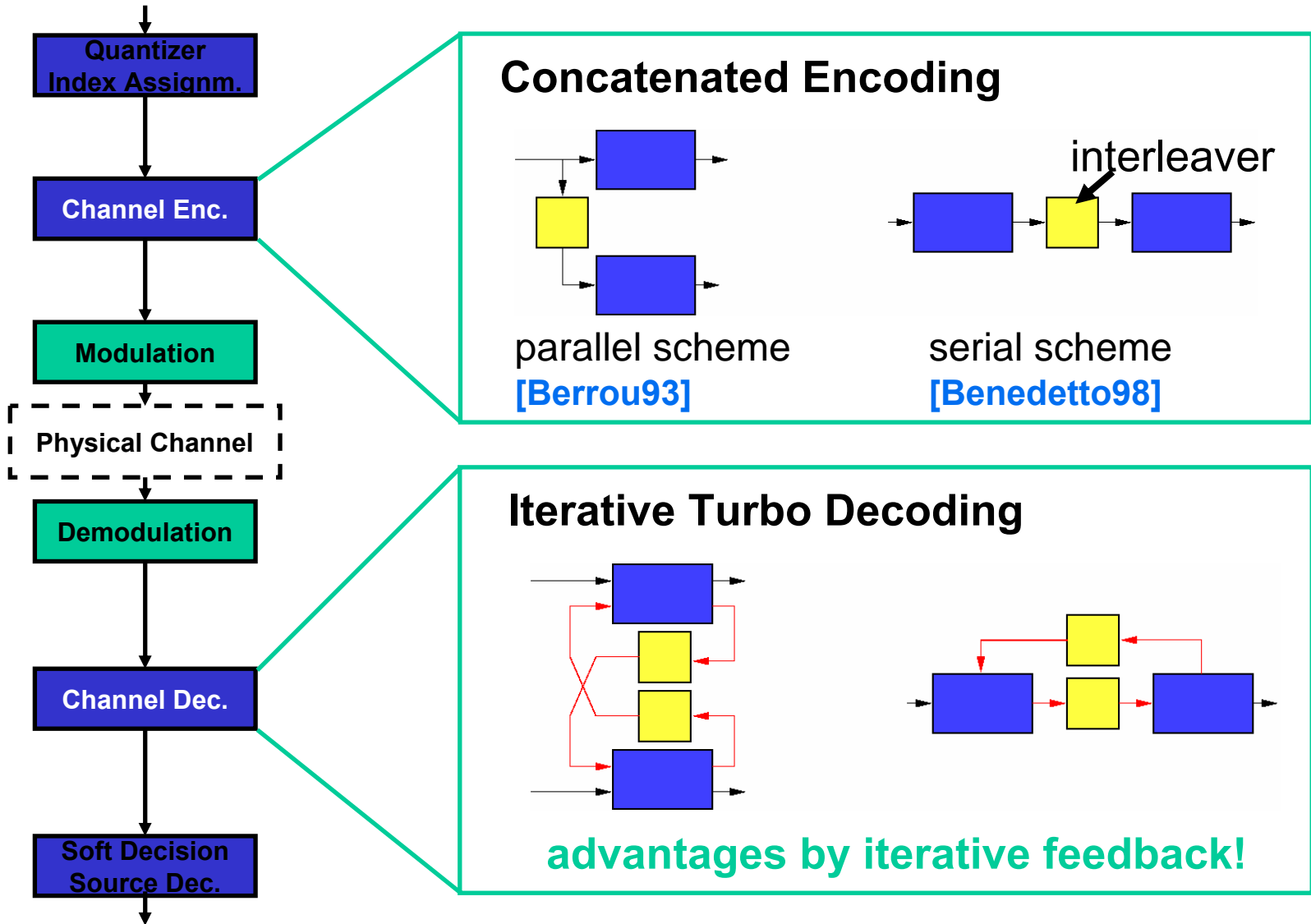
“Information on a bit originating from other bit positions”

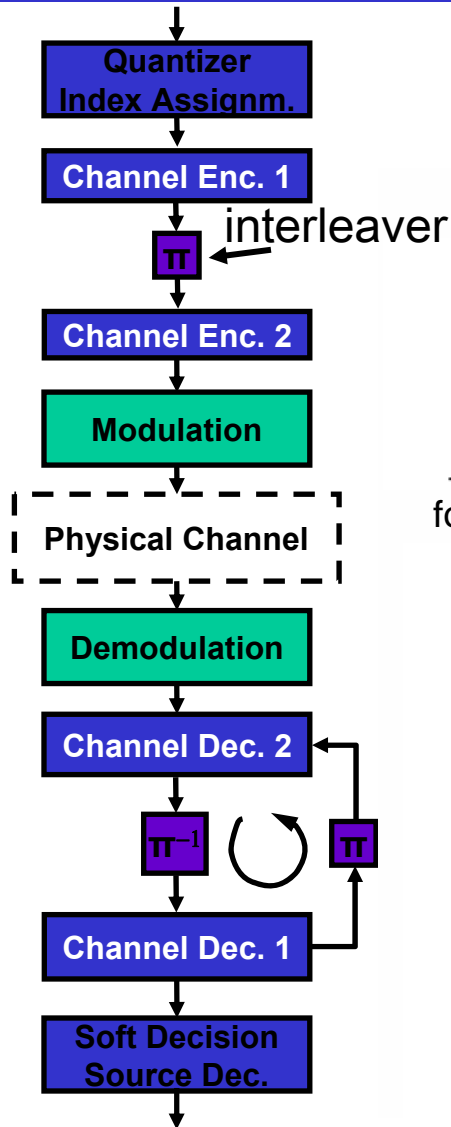
- Types of a priori information



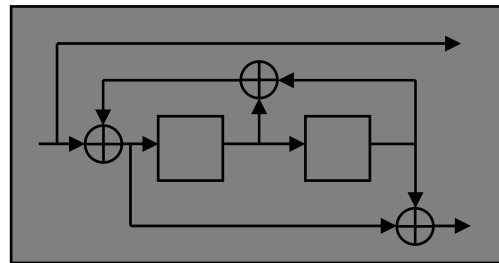
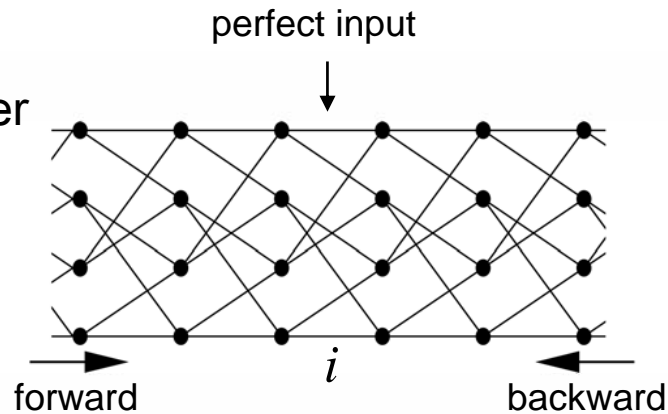
$$L(\hat{x}_i|z_i) = L(z_i|x_i) + L(x_i) + L_i^{[ext]}$$

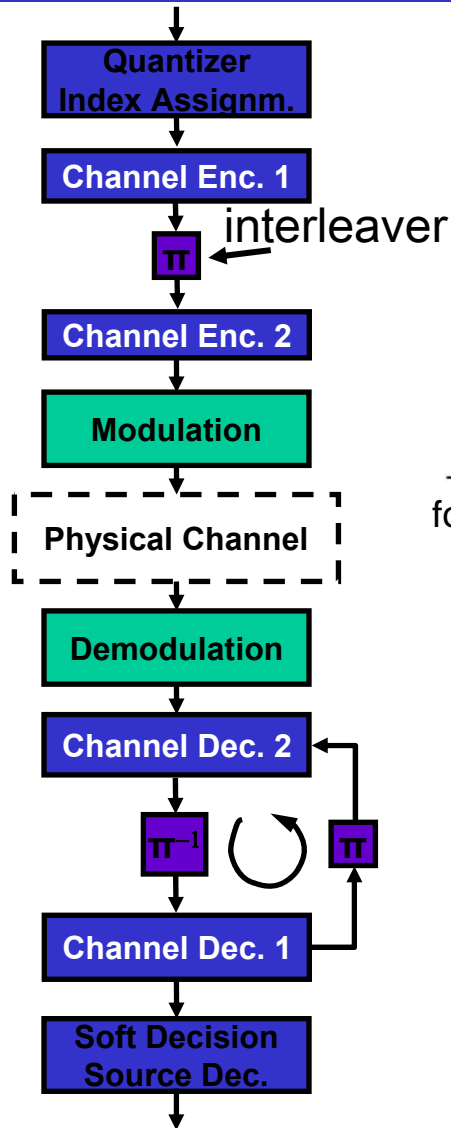
- Gained information from other bit positions through processing in component
- Dynamic a priori information is the new key element
 - Other Turbo component's updated and improved extrinsic information
- Independent (i.e. uncorrelated) extrinsic information assumed and required



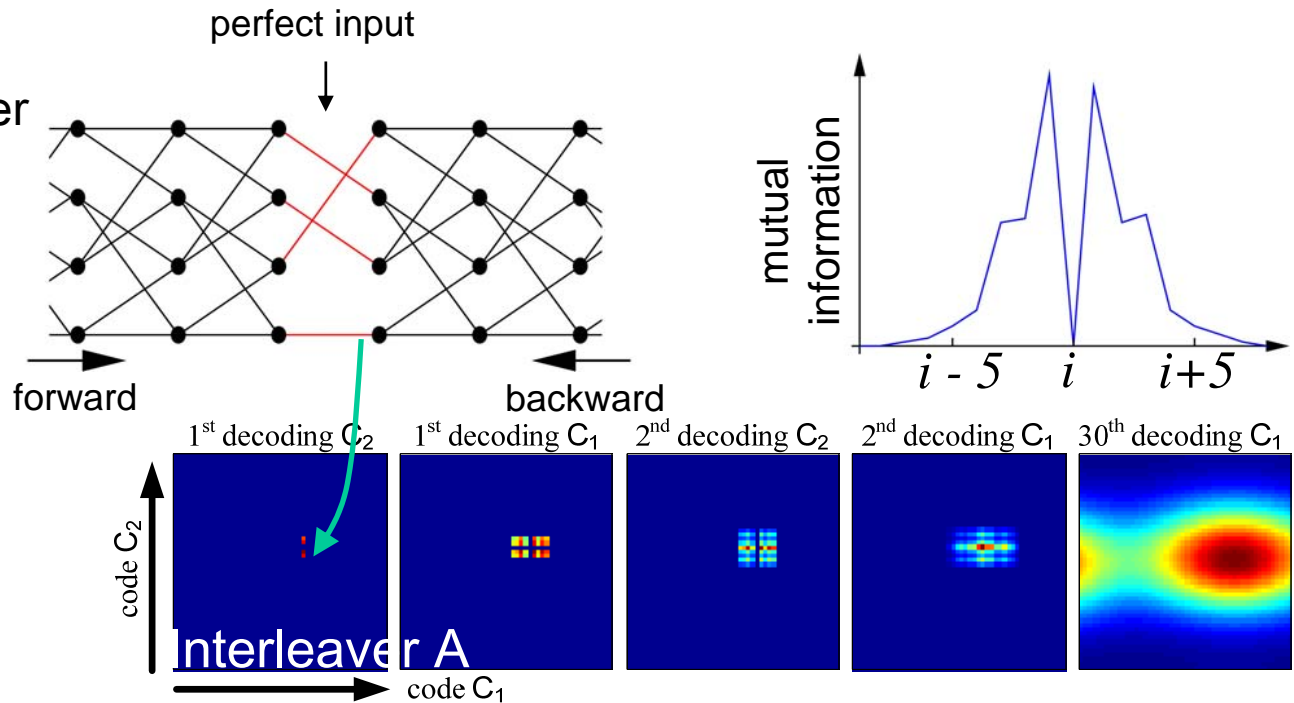


Propagation of extrinsic information

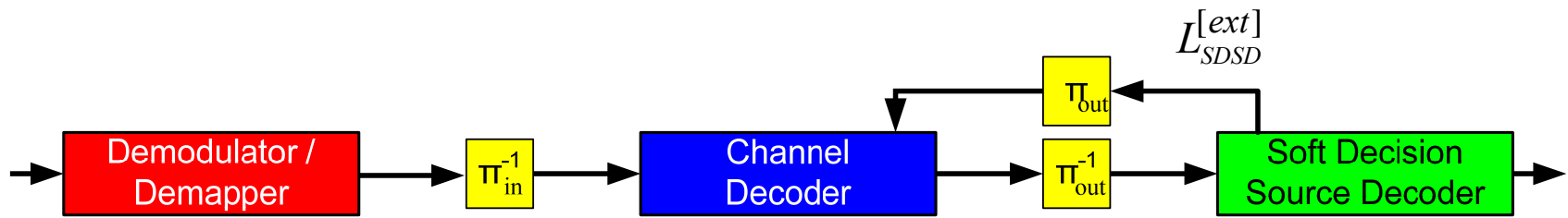




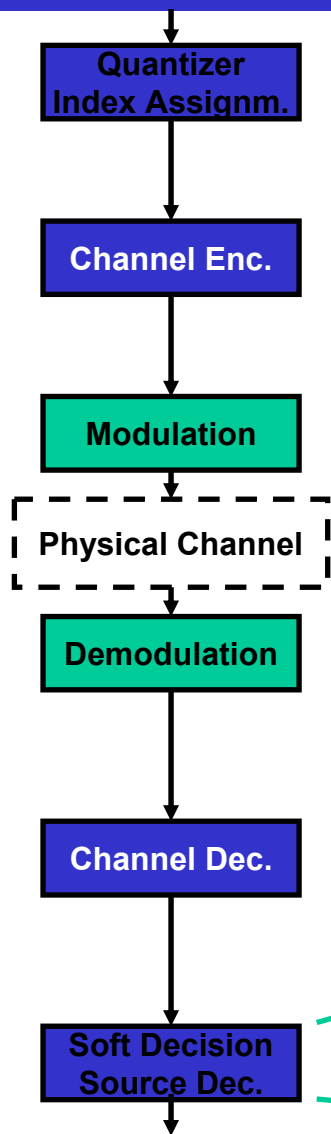
Propagation of extrinsic information



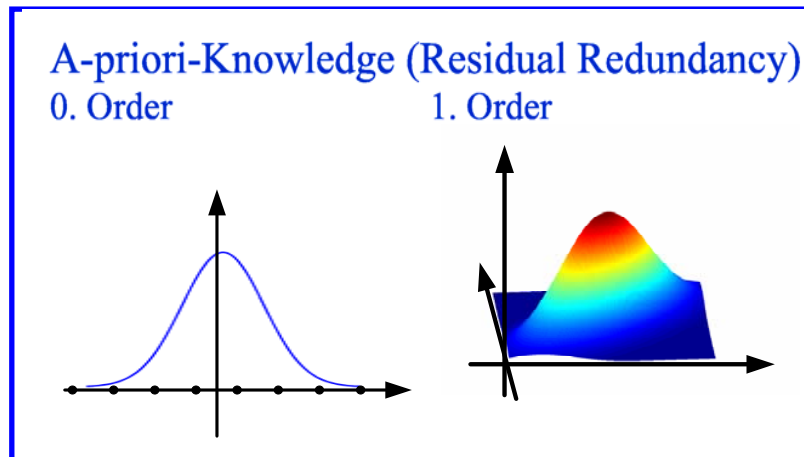
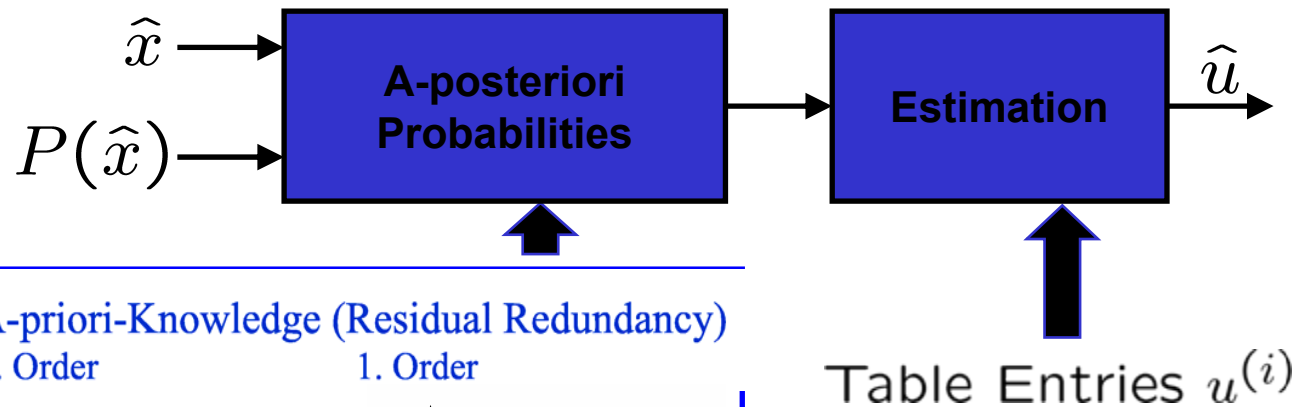
Iterative Source-Channel Decoding (ISCD)



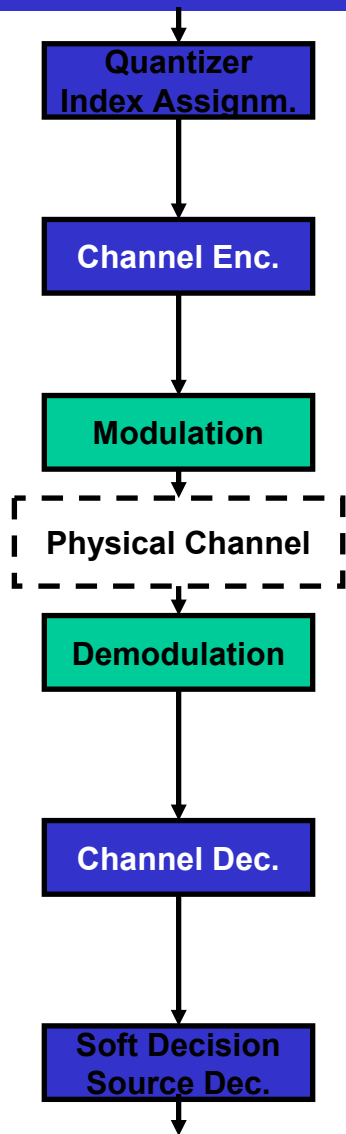
- joint
- channel decoding
 - source decoding



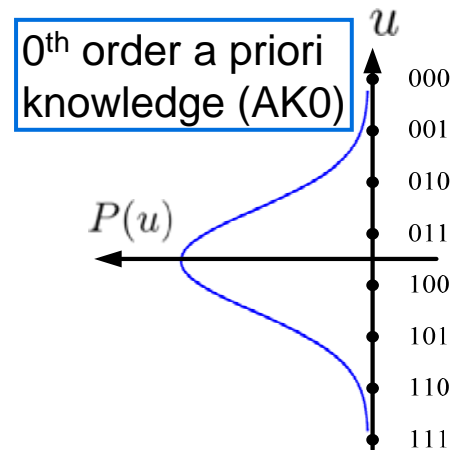
- Parameter Decoding by Estimation
 - Minimum Mean Square Error (MMSE) criterion



$$\hat{u} = \sum_i u^{(i)} \cdot P\left(x_t^{(i)} \mid \hat{x}_t, \hat{X}_{t-1}\right)$$

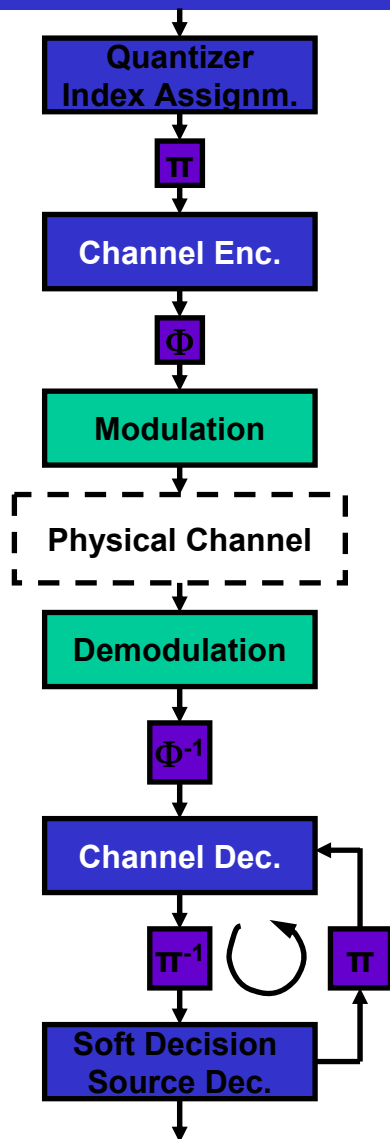


- Exploits residual redundancy in source samples
- Based on fully connected Trellis diagram

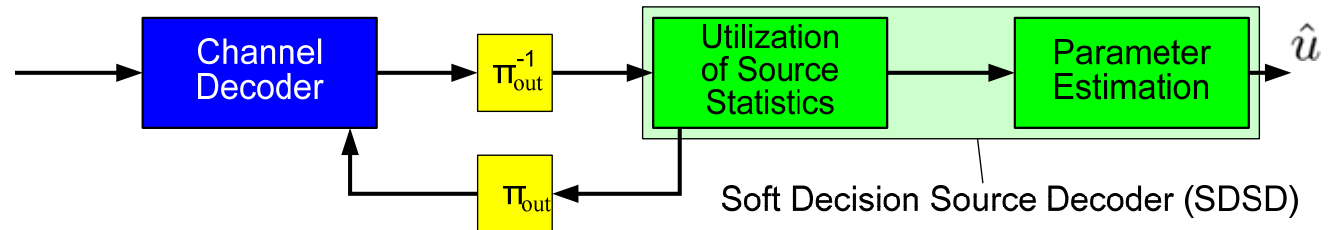


- Soft parameter estimation instead of table lookup (minimizes mean-square error)

$$\hat{u} = \sum_{\xi=0}^{Q-1} \bar{u}^{(\xi)} \cdot P(\bar{u}^{(\xi)} | \underline{z})$$

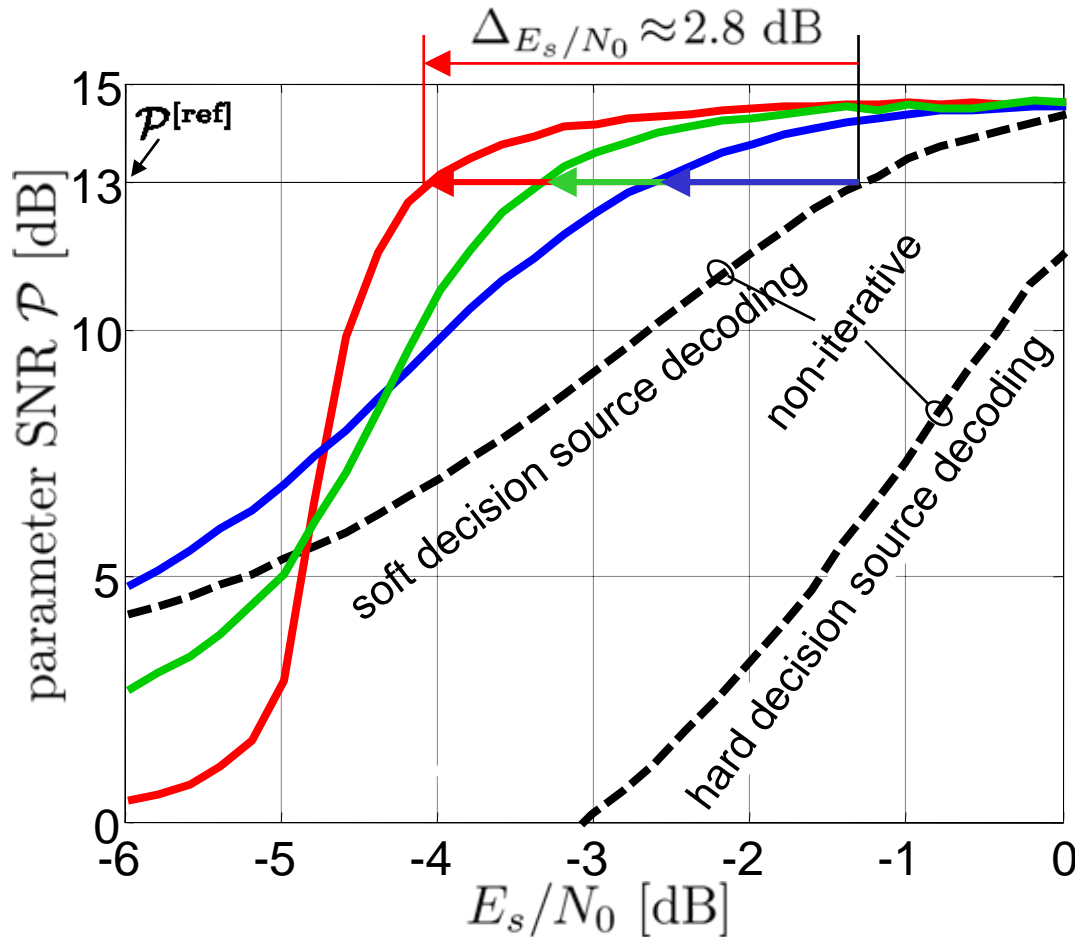


- Iterative evaluation of channel decoding and soft decision source decoding



- Key elements determining the performance of ISCD
 - Index assignment of bit patterns to quantization levels
 - Preserve (exploitable) residual redundancy in source samples after source encoding

- Evaluation at reference parameter SNR $\mathcal{P}^{[\text{ref}]}$

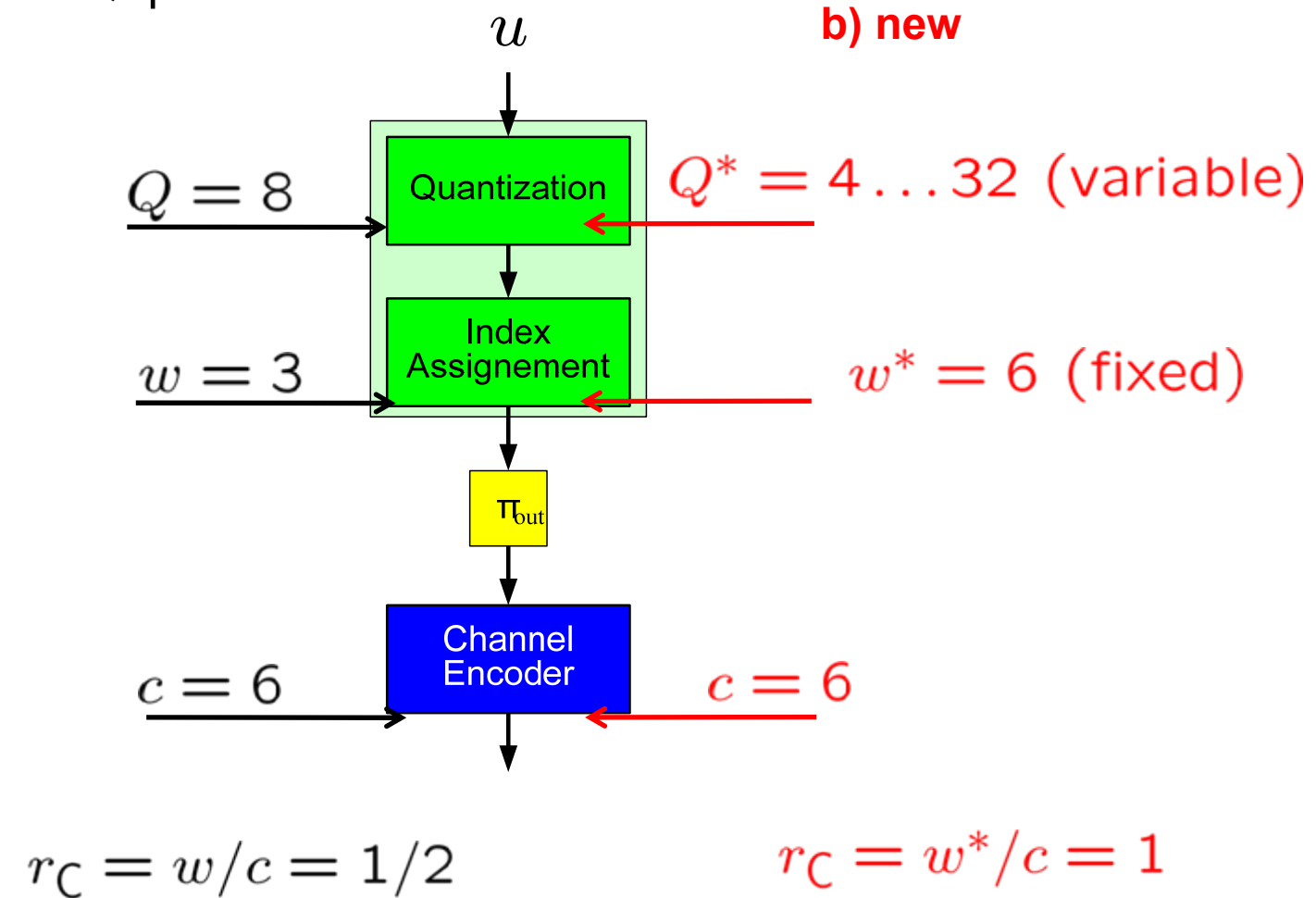
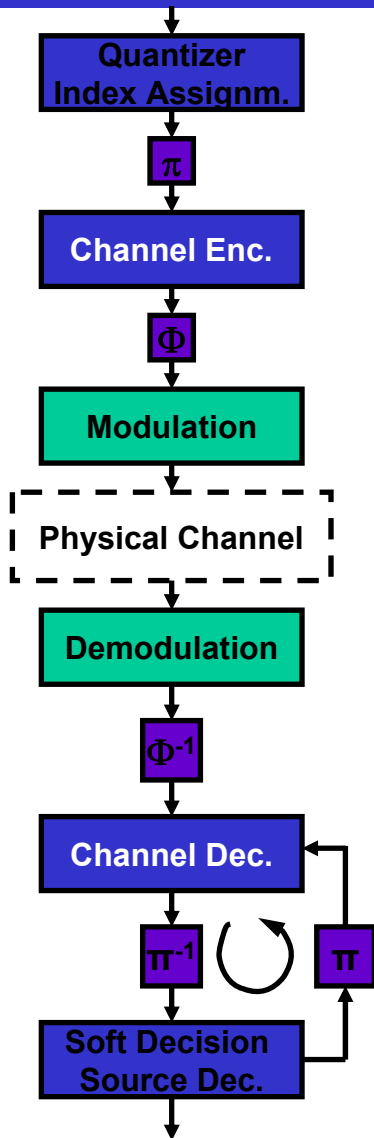


- AWGN channel
- BPSK modulation
- 250 parameters/frame
- auto-correlation $\rho=0.9$
- 3 bit Lloyd-Max quant. (LMQ)
- $r_c=1/2$ conv. code with 8 states
- 10 iterations
- **apply ISCD**
 - rec. syst. conv. (RSC) code
 - natural binary (NB) index assignm.
- **improved index assignment**
 - rec. syst. conv. (RSC) code
 - EXIT optim. (EO) index assignm.
- **improved channel code**
 - rec. **non-syst.** conv. (RNSC) code
 - EXIT optim. (EO) index assignm.

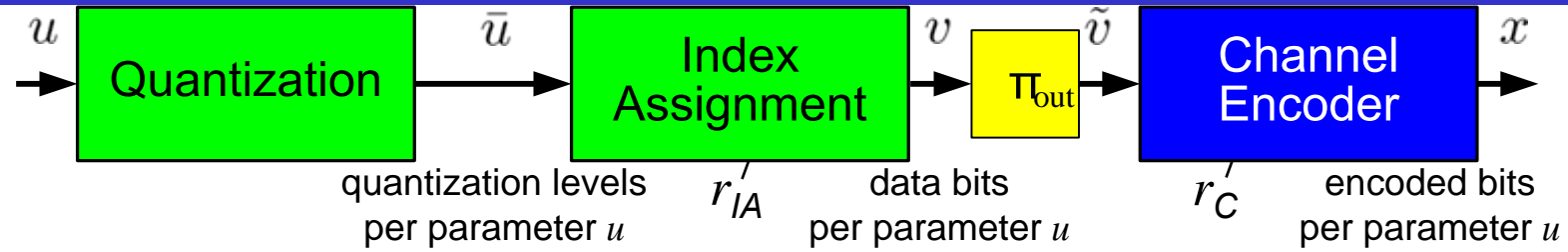
Index assignment and channel coding are key factors for ISCD performance

FlexCode Redundant Index Assignments [Adrat05], [Clevorn06]

- Highly Redundant Index Assignments
- Q quantization levels



FlexCode ISCD - Index Assignments for Multi-Mode Extension



- classic ISCD ($r_C=1/2$) $Q = 8$ $w = \text{ld}(Q) = 3$ $w/r_C = 6$
- optimized ISCD ($r_C=1$) $Q^* = 8$ $w^* = w/r_{IA} = 6$ $w^*/r_C = 6$
- multi-mode ISCD ($r_C=1$) $Q^* = \text{variable}$ $w^* = w/r_{IA} = 6$ $w^*/r_C = 6$

Redundant index assignment optimized, e.g., block coded

$Q^* \leq 2$
multi-mode, e.g., block coded with $Q^* = 16$

ξ	$\{\xi\}_2$	v
0	000	000000
1	001	001011
\vdots	\vdots	\vdots
7	111	111000

$$\mathbf{G}_{\text{BC}(6,3)} = \begin{pmatrix} 100101 \\ 010110 \\ 001011 \end{pmatrix}$$

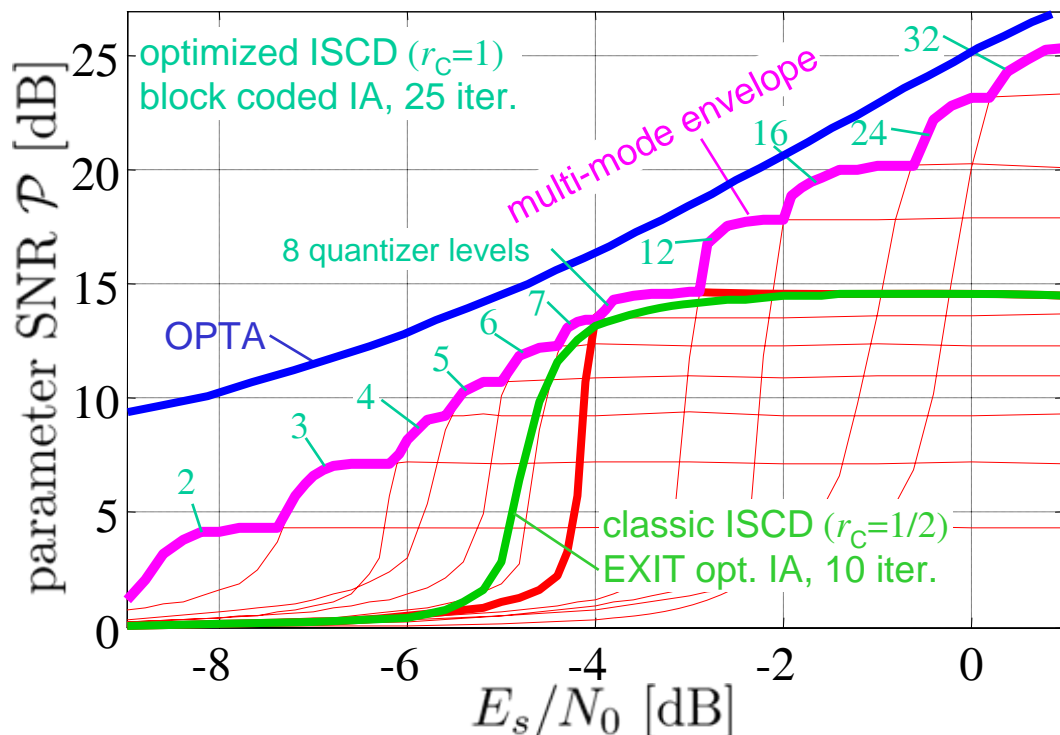
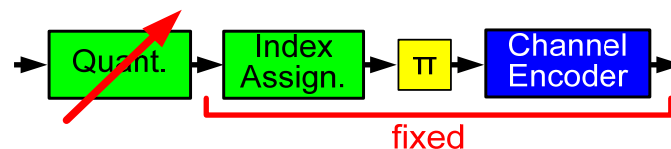
ξ	$\{\xi\}_2$	v
0	0000	000000
1	0001	001011
\vdots	\vdots	\vdots
7	0111	111000
\vdots	\vdots	\vdots
11	1011	100010
\vdots	\vdots	\vdots
15	1111	111000

with $Q^* = 12$

$Q^* = 8$ with $\mathbf{G}_{\text{BC}(6,3)}$ included

$$\mathbf{G}_{\text{BC}(6,4)} = \begin{pmatrix} 1111111 \\ 100101 \\ 010110 \\ 001011 \end{pmatrix}$$

- With block coded index assignment only single lookup table required
- Interleaver and channel code are fixed
- Only quantizer levels must be adapted



- AWGN channel
- BPSK modulation
- 250 parameters/frame
- auto-correlation $\rho=0.9$
- Q level Lloyd-Max quant. (LMQ)
- 6 bits per parameter
- convolutional codes with 8 states

- Higher quality in good channel conditions due to large number of quantizer levels
- More robust transmission in bad channel conditions due to increased redundancy

- Iterative Source-Channel Decoding (ISCD) permits **near-capacity** decoding for BPSK modulation
- Redundant index assignments allow versatile multi-mode extension
- Additional flexibility due to multi-mode ISCD
 - good channels: small quantization error
 - bad channels: large quantization errors, but insensitive to channel noise

THANK YOU!